

## Discrete mathematics I. – Homework 3

Problem sessions in the week of October 14, 2013

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**1.** Let  $T$  be a satisfiable set of formulae. Show that  $T \cup \{A \vee B\} \models C$  if and only if simultaneously  $T \cup \{A\} \models C$  and  $T \cup \{B\} \models C$ .

**2.** Let  $T$  be a satisfiable set of formulae. Show that if the set of formulae  $T \cup \{\neg A\}$  is unsatisfiable then  $T \models A$ . (the other implication was done during the lecture)

**3.** Show that the formula  $\neg a$  can be inferred from the premisses  $\{\neg(a \Rightarrow b), \neg(a \Rightarrow b) \Rightarrow \neg a\}$ , i.e. find a formal proof according to the definition.

**4.** Let  $D$  be any axiom of the propositional calculus and  $E$  is any formula. Find a proof of the formula  $E \Rightarrow D$ .

**5.** Let  $T$  be a set of formulae and  $A, B, C$  are any formulae. Show that if  $T \vdash A \Rightarrow B$  and  $T \vdash A \Rightarrow (B \Rightarrow C)$  then  $T \vdash A \Rightarrow C$ . Hint: use axiom (A2).

**6.** Show that for the formula  $(A \Rightarrow B) \Rightarrow (\neg A \Rightarrow \neg B)$  it is not possible to find any proof in the formal system, i.e. it is not a theorem of propositional calculus. Hint: has to be any formula in the sequence of the proof tautology?

**7.** Show that the formula  $a \Rightarrow b$  can be deduced from premise  $\{\neg(a \Rightarrow b) \Rightarrow \neg a\}$ . Hint: show  $\{\neg(a \Rightarrow b) \Rightarrow \neg a\} \vdash a \Rightarrow (a \Rightarrow b)$ , use axiom (A2) and problem 3.

Compare with problem 7a) from the last homework.